

CUR 412: Game Theory and its Applications, Lecture 14

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June 19, 2015

Announcements

- ▶ Homework #5 is due on Tuesday, June 23. You may put it in my mailbox on the 9th floor of Boxue, or in the envelope on my office door at 123 Qiusuo.
- ▶ The final will be on July 2, 8:00-10:00 AM, in Boxue 604.
- ▶ Final will be closed-book and covers the second half of the course. Old finals and solutions are on the course website.
- ▶ Bring a non-programmable calculator.

Review of Last Lecture

- ▶ A *repeated game* is a situation where a game is repeatedly played for several (possibly infinite) periods.
- ▶ We assume that agents value payoffs in the *present*, more than payoffs in the *future* by a constant *discount factor*, δ .
- ▶ δ is between 0 and 1; an agent with a lower δ is said to be more *impatient*, i.e. places lesser weight on the future.
- ▶ The overall value of a sequence of payoffs w^1, w^2, w^3, \dots is given by the *discounted average*:

$$(1 - \delta) \sum_{t=1}^{\infty} \delta^{t-1} w^t$$

- ▶ A strategy for a repeated game must specify an action for every possible history.

Review of Last Lecture

- ▶ We saw some examples of strategies: *Grim Trigger*, *Punish for k periods*, *Tit-for-Tat*.
- ▶ Strategies can be defined as a function mapping any history to an action.
- ▶ Or, we can represent them with a state diagram.
- ▶ To check whether a pair of strategies forms a NE in the infinitely repeated Prisoner's Dilemma, we must:
 - ▶ Calculate the discounted average of the payoff stream to both players when they *do not* deviate;
 - ▶ Show that there is no alternative strategy for either player that gives a higher discounted average.

Review of Last Lecture

- ▶ The value of δ , the discount factor, can affect whether a strategy pair can be a NE or not.
- ▶ Generally, δ must be higher than some level for a NE to exist.
- ▶ That is, the threat of future punishment can deter defection in the present only if players place a high enough weight on the future.
- ▶ Some known NE: when both players use *Grim Trigger* or *Tit-for-Tat*.

Ch 14.9: Subgame Perfect Equilibria and the One-Deviation Property

- ▶ We know that Nash equilibria of an extensive form game may include threats that are not credible.
- ▶ This is particularly important in repeated games, since the threat of punishment is the only way to deter a player from playing *Defect*.
- ▶ If punishment is not credible, then there will be no reason to expect anything but (D, D) .
- ▶ The concept of subgame perfect equilibrium dealt with this issue by requiring that strategies also be Nash equilibria in all subgames.
- ▶ However, this is not easy to check in an infinite game.
- ▶ The good news is that we can take advantage of a result that gives a much simpler condition that is easier to check.

The One-Deviation Property

- ▶ Suppose we have a strategy profile s . s satisfies the **one-deviation property** if no player can increase his payoff in any subgame through a *one-shot deviation*:
 - ▶ deviating from s in the *first* period of the subgame, then *reverting* back to his strategy in s for the rest of the game.
- ▶ In a *finite* horizon game, a strategy profile is a SPNE if and only if it satisfies the one-deviation property.
- ▶ In an *infinite* horizon game where the discount factor is less than 1, a strategy profile is a SPNE if and only if it satisfies the one-deviation property.
- ▶ This is also called the *one-shot deviation principle*.

The One-Deviation Property

- ▶ The basic idea is that any change in strategy can be broken down into a sequence of one-period changes.
- ▶ If there is no one-period change that can increase payoffs, then there is no change in strategy that can increase payoffs.
- ▶ A full proof is somewhat technically advanced; however, if you are interested, you can search for "one-shot deviation principle".
- ▶ For repeated games, this means that in order to check whether some strategy profile s is a SPNE, we must check all possible one-shot deviations for all players.
- ▶ for every possible history, compare the payoffs to adhering to s , versus deviating for one period, then reverting back to s .
- ▶ Let's look at some of the strategies we've seen, and check if it is a SPNE when both players play them.

Grim Trigger

- ▶ *Grim Trigger* is a strategy that punishes defection, so it seems like it should be possible to be part of a SPNE.
- ▶ We will show that it is *not* a SPNE when both players play Grim Trigger.
- ▶ Suppose the previous history ended with (C, D) .
- ▶ We don't specify how this occurred; in fact, if both players adhere to Grim Trigger, then the outcome will be (C, C) in every period. However, both players always have the choice of playing D , so it is one of the possible subgames.

Grim Trigger

- ▶ First, calculate the payoffs from *not* deviating. If both players adhere to *Grim Trigger*:
 - ▶ outcome path will be: $(D, C), (D, D), (D, D), \dots$ (Player 1 will start punishing in the first period, Player 2 will start in the second period)
 - ▶ payoffs will be: $(3, 0), (1, 1), (1, 1), \dots$
 - ▶ Player 1's discounted average:

$$(1 - \delta)(3 + \delta + \delta^2 + \dots) = (1 - \delta)\left(3 + \frac{\delta}{1 - \delta}\right) = 3 - 2\delta$$

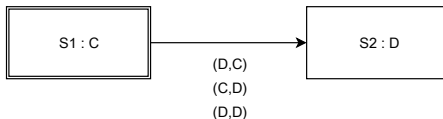
- ▶ Player 2's discounted average:

$$(1 - \delta)(0 + \delta + \delta^2 + \dots) = \delta$$

Grim Trigger

- ▶ Now, suppose Player 2 deviates by choosing D in all periods (note that this is not a one-shot deviation).
 - ▶ outcome path: $(D, D), (D, D), (D, D), \dots$
 - ▶ payoffs: $(1, 1), (1, 1), \dots$
 - ▶ Both players' discounted average is 1.
- ▶ By our assumptions on $\delta, 1 > \delta$.
- ▶ Therefore, Player 2 has an incentive to deviate, given a history ending with (C, D) , so this violates the requirement for a SPNE.

Modified Grim Trigger



- ▶ However, there is a modified version of Grim Trigger that is a SPNE when both players play it.
- ▶ *Grim Trigger* moves into the punishment state if the *other* player plays *D*.
- ▶ In contrast, *Modified Grim Trigger* moves into the punishment state if *either* player plays *D*.

Modified Grim Trigger

- ▶ Now, let's use the One-Deviation Property to show that it is a SPNE when both players play *Modified Grim Trigger*.
- ▶ We need to check all possible possible histories for the case where a one-shot deviation increases a player's payoff.
- ▶ This is vastly simplified because the strategy distinguishes between two types of histories:
 - ▶ D has never been played by either player (therefore, play C)
 - ▶ D has been played by either player (therefore, play D)
- ▶ For each of these types of histories, we will check the conditions that satisfy the one-deviation principle.

Case 1: D has not been played

- ▶ Case 1: Suppose D has never been played by either player.
- ▶ This is the case if the history is empty (i.e. at the beginning of the game), or if only (C, C) has been played in all periods.
- ▶ First, let's calculate the payoffs from *not* deviating.
 - ▶ outcome path: $(C, C), (C, C), \dots$
 - ▶ payoffs: $(2, 2), (2, 2), \dots$
 - ▶ Both players' discounted average is 2.

Case 1: D has not been played

- ▶ Now, consider the payoffs from a one-shot deviation. Suppose Player 1 does the one-shot deviation and plays D in the first period, then reverts back to *Modified Grim Trigger*.
 - ▶ outcome path: $(D, C), (D, D), (D, D), \dots$
 - ▶ payoffs: $(3, 0), (1, 1), (1, 1), \dots$
 - ▶ Player 1's discounted average is $3 - 2\delta$, same as shown above for Grim Trigger.
- ▶ Therefore, this one-shot deviation will give a higher payoff if $3 - 2\delta > 2$, or if $\delta < \frac{1}{2}$.
- ▶ The same is true if Player 2 does the one-shot deviation.

Case 2: D has been played

- ▶ Case 2: D has been played in the past, by either player.
- ▶ Then, *Modified Grim Trigger* will switch to the punishment state. If the players do *not* deviate:
 - ▶ outcome path: $(D, D), (D, D), \dots$
 - ▶ payoffs: $(1, 1), (1, 1), \dots$
 - ▶ Both players' discounted average is 1.

Case 2: D has been played

- ▶ Suppose Player 1 does a one-shot deviation by playing C .
 - ▶ outcome path: $(C, D), (D, D), (D, D), \dots$
 - ▶ payoffs: $(0, 3), (1, 1), (1, 1), \dots$
 - ▶ Player 1's discounted average is $(1 - \delta)(\delta + \delta^2 + \dots) = \delta$
- ▶ This one-shot deviation will not give a higher payoff for any value of δ .
- ▶ The same is true if it is Player 2 that considers a one-shot deviation.

Modified Grim Trigger

- ▶ The requirements for the two cases we've seen are:
 - ▶ Case 1: D has never been played. A one-shot deviation is profitable if $\delta < \frac{1}{2}$.
 - ▶ Case 2: D has been played in the past, by either player. A one-shot deviation is never profitable.
- ▶ Combining the requirements for both cases, a profitable (i.e. higher payoff) one-shot deviation does not exist if $\delta \geq \frac{1}{2}$, and therefore the strategy profile is a SPNE.

Tit for Tat

- ▶ Recall: *Tit – for – Tat* plays C at the beginning of the game, then plays whatever the other player did in the previous period.
- ▶ Therefore, in terms of determining the outcome, we need only look at the previous round's actions. There are four possibilities:
 1. History ends in (C, D)
 2. History ends in (D, C)
 3. History ends in (C, C)
 4. History ends in (D, D)
- ▶ Let's examine each case in turn. We will find the conditions on δ that ensure a profitable one-shot deviation does not exist.
- ▶ We need to check both players' possible one-shot deviations.

Case 1: Previous outcome was (C, D)

- ▶ Suppose the history ends with (C, D) .
- ▶ If both players do *not* deviate:
 - ▶ outcome path: $(D, C), (C, D), (D, C), \dots$ (both players will alternate between C and D forever)
 - ▶ payoffs: $(3, 0), (0, 3), (3, 0), \dots$
 - ▶ Player 1's discounted average:

$$(1 - \delta)(3 + 0 + 3\delta^2 + 0 + 3\delta^4 + \dots) = (1 - \delta) \frac{3}{1 - \delta^2} = \frac{3}{1 + \delta}$$

- ▶ Player 2's discounted average:

$$(1 - \delta)(0 + 3\delta + 0 + 3\delta^3 + 0 + \dots) = (1 - \delta) \frac{3\delta}{1 - \delta^2} = \frac{3\delta}{1 + \delta}$$

Case 1: Previous outcome was (C, D)

- ▶ Suppose Player 1 does a one-shot deviation by playing C in the first period.
 - ▶ outcome path: $(C, C), (C, C), \dots$
 - ▶ payoffs: $(2, 2), (2, 2), \dots$
 - ▶ Player 1's discounted average is 2.
- ▶ Therefore, this one-shot deviation will give a higher payoff if $2 > \frac{3}{1+\delta}$, or if $\delta > \frac{1}{2}$.
- ▶ Now, suppose Player 2 does a one-shot deviation by playing D in the first period:
 - ▶ outcome path: $(D, D), (D, D), \dots$
 - ▶ payoffs: $(1, 1), (1, 1), \dots$
 - ▶ Player 2's discounted average is 1.
- ▶ Therefore, this one-shot deviation will give a higher payoff if $1 > \frac{3\delta}{1+\delta}$, or if $\delta < \frac{1}{2}$.
- ▶ Combining the two conditions, a profitable one-shot deviation does not exist if $\delta = \frac{1}{2}$.

Case 2: Previous outcome was (D, C)

- ▶ Suppose the previous outcome was (D, C) .
- ▶ This is the same situation as the previous case, with the players reversed.
- ▶ Therefore, the result is the same: a profitable one-shot deviation does not exist if $\delta = \frac{1}{2}$.

Case 3: Previous outcome was (C, C)

- ▶ Suppose the previous outcome was (C, C) .
- ▶ If both players do *not* deviate:
 - ▶ outcome path: $(C, C), (C, C), \dots$
 - ▶ payoffs: $(2, 2), (2, 2), (2, 2), \dots$
 - ▶ Both players' discounted average: 2

Case 3: Previous outcome was (C, C)

- ▶ Suppose Player 1 does a one-shot deviation by playing D in the first round.
 - ▶ outcome path: $(D, C), (C, D), (D, C), (C, D), \dots$
 - ▶ payoffs: $(3, 0), (0, 3), (3, 0), \dots$
 - ▶ Player 1's discounted average: $\frac{3}{1+\delta}$ (shown in the first type of history).
- ▶ The one-shot deviation will be profitable if $\frac{3}{1+\delta} > 2$, or if $\delta < \frac{1}{2}$.
- ▶ The same is true if it is Player 2 that is considering a one-shot deviation.
- ▶ Therefore, a profitable one-shot deviation does not exist if $\delta \geq \frac{1}{2}$.

Case 4: Previous outcome was (D, D)

- ▶ Suppose the previous outcome was (D, D) .
- ▶ If both players do *not* deviate:
 - ▶ outcome path: $(D, D), (D, D), \dots$
 - ▶ payoffs: $(1, 1), (1, 1), (1, 1), \dots$
 - ▶ Both players' discounted average: 1

Case 4: Previous outcome was (D, D)

- ▶ Suppose Player 1 does a one-shot deviation by playing C in the first round.
 - ▶ outcome path: $(C, D), (D, C), (C, D), \dots$
 - ▶ payoffs: $(0, 3), (3, 0), (0, 3), (3, 0), \dots$
 - ▶ Player 1's discounted average: $\frac{3\delta}{1+\delta}$ (see Case 1 for the formula for the discounted avg)
- ▶ The one-shot deviation will be profitable if $\frac{3\delta}{1+\delta} > 1$, or if $\delta > \frac{1}{2}$.
- ▶ The same is true if it is Player 2 that is considering a one-shot deviation.
- ▶ A profitable one-shot deviation will not exist if $\delta \leq \frac{1}{2}$.

- ▶ Summarizing the results for each case:
 - ▶ Case 1 (previous outcome was (C, D)): no profitable deviation exists if $\delta = \frac{1}{2}$.
 - ▶ Case 2 (previous outcome was (D, C)): no profitable deviation exists if $\delta = \frac{1}{2}$.
 - ▶ Case 3 (previous outcome was (C, C)): no profitable deviation exists if $\delta \geq \frac{1}{2}$.
 - ▶ Case 4 (previous outcome was (D, D)): no profitable deviation exists if $\delta \leq \frac{1}{2}$.
- ▶ Combining all the conditions, no profitable deviation exists if and only if $\delta = \frac{1}{2}$.
- ▶ This result depends on the particular payoffs of the game in each period.
- ▶ If we change the payoffs, the range of δ that supports a SPNE may change.

Tacit Collusion Among Firms

- ▶ One important application of repeated-game equilibria is the study of *tacit collusion* among firms.
- ▶ Usually, it is considered illegal for competitors to fix prices (i.e. agree not to compete on price, and keep prices high).
- ▶ If executives from rival firms are seen to meet and discuss prices, this is grounds for a lawsuit.
- ▶ However, even if they do not communicate, firms may collude, simply by observing the past history of prices.
- ▶ "Price wars" can be seen as times when firms seek to punish each other.

- ▶ Consider this Cournot duopoly problem with two firms.
- ▶ Each firm has cost of production $c_i(q_i) = 10q_i$.
- ▶ Market demand is given by $P = 100 - Q$.
- ▶ The two firms repeatedly play the Cournot duopoly game in time periods $t = 1, 2, 3, \dots$ with discount factor δ .
- ▶ The NE of the Cournot game in a single period is $q_1 = q_2 = 30$, $p = 40$, and profits for each firm are 900.
- ▶ If there was a single monopolist firm, the optimal $q = 45$, $p = 55$, profits = 2025.
- ▶ Suppose in each period, each firm can choose to *Collude*, in which case they produce $q_i = 22.5$, half the monopoly quantity.
- ▶ Or, they can choose to *Defect*, in which case they maximize their own profit, given the quantity of the other firm.

- ▶ If both firms choose *Defect*, they choose the NE quantities ($q_i = 30$), which results in a profit of 900 for each firm.
- ▶ If both firms choose *Collude*, they choose half the monopoly quantities ($q_i = 22.5$), which results in a profit of 1012.5 for each firm.
- ▶ If one chooses *Collude* while the other chooses *Defect*, the firm that *Colludes* chooses $q_i = 22.5$, while the firm that *Defects* chooses the best response to that, which is $q_j = 33.75$.
- ▶ The firm that chose *Collude* gets a profit of 759.375, while the firm that chose *Defect* gets a profit of 1139.06.
- ▶ Note that this situation is a Prisoner's Dilemma.
- ▶ We can then use the one-shot deviation principle to show that a given strategy profile (e.g. both players play Modified Grim Trigger) is a SPNE, for a certain value of the discount factor δ .

- ▶ The one-shot deviation principle is not limited to repeated Prisoner's Dilemma; it can be used in any repeated game situation to prove that a given strategy profile is or is not a SPNE.

A Review of Some Key Ideas

- ▶ Recall the original goal of the course: we want to be able to predict the outcomes of strategic situations.
- ▶ We have seen a variety of techniques and ideas that are used in order to try to answer this question.
- ▶ Let's review some of the topics we've covered and see some of the ideas behind them.

Idea 1: Play the Optimal Action

- ▶ The simplest idea is to simply see if there is an action that is always the optimal choice in any situation.
- ▶ If it exists, a rational player will always play it.
- ▶ This is a *strictly dominant* action.
- ▶ If all players have a strictly dominant action, then we can predict that will be the outcome.
- ▶ However, most games don't have a strictly dominant action.

Idea 2: Eliminate Sub-optimal Actions

- ▶ If we can't find a clearly optimal action, we can try to get rid of actions that are clearly sub-optimal.
- ▶ A *strictly dominated* action is never a best response to any situation, so a rational player will never play it.
- ▶ Therefore, it will not be part of any kind of equilibrium (NE, SPNE, etc).
- ▶ If we make the assumption that players know that other players are rational, then we can repeatedly eliminate dominated actions.
- ▶ The set of outcomes that remain are said to be *rationalizable*.

Idea 3: Find an Equilibrium of the System

- ▶ What if we are left with multiple outcomes even after eliminating all the dominated actions? (e.g. BoS)
- ▶ Instead of trying to directly find the "best" outcome, we can try a less ambitious goal: find a *stable point*, or equilibrium, of the system.
- ▶ An equilibrium is a point such that the system stays there, once it reaches that point.
- ▶ For example, when analyzing a market with a supply and demand curve, the equilibrium is the intersection.
- ▶ If the price diverges, there will be excess demand or supply, driving the price back to equilibrium.

Idea 3: Find an Equilibrium of the System

- ▶ Nash equilibrium is a specific type of equilibrium where we define "moving away" as due to the action of a *single* player.
- ▶ We can imagine different conditions for equilibrium; for example, if we consider the possibility of multiple players cooperating to change the outcome.
- ▶ This is the subject of *coalitional game theory*.
- ▶ Note that this sidesteps many important questions:
 - ▶ How do players discover the equilibrium outcomes?
 - ▶ How long does it take to reach equilibrium?
 - ▶ What if there are many equilibria; which ones are "better" than others?

Idea 4: Allow Combinations of Actions

- ▶ A mixed strategy is a way of combining actions, by creating a weighted average of actions.
- ▶ We have seen that a mixed strategy can dominate pure strategies, even if the pure strategies don't have a dominant action.
- ▶ Allowing combinations of actions also guarantees the existence of an equilibrium: this is Nash's proof that a NE exists in a finite game with mixed strategies.

Idea 5: Predict What Happens in Any Situation

- ▶ In extensive-form games, we use backwards induction to find a solution.
- ▶ This is equivalent to predicting the outcome for all possible situations (e.g. subgames), and assuming that this outcome will hold.
- ▶ The subgame perfect concept requires that equilibrium conditions must hold at every step of the game.
- ▶ However, as we've seen in the Centipede Game, it becomes harder to predict what happens many steps in the future.

Idea 6: Beliefs as Probabilities

- ▶ When there is uncertainty (about other players, or about the environment), one way to handle this is to use probability distributions to model beliefs.
- ▶ Bayes' Rule gives us a way to calculate the correct probabilities of any event, given knowledge of other events that may influence the outcome.
- ▶ We can expand our definition of equilibrium to include beliefs as well as strategies; this gives us our concept of weak sequential equilibrium.

Idea 7: There's Always a Future

- ▶ In the finitely repeated Prisoner's Dilemma, we've seen that the only NE and SPNE result in (D, D) every period.
- ▶ It is only possible to deter defection and maintain cooperation in the *infinitely* repeated game, because there is always the possibility of future punishment.
- ▶ This also requires that players have a sufficiently high discount factor; otherwise, the threat of future punishment is outweighed by the prospect of immediate gain.

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